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CRUK CI HPC cluster introduction (III of III)

Some advanced topics



Reference Genomes

• Path to reference data:

/scratchb/bioinformatics/reference_data/reference_genomes/

• Path to assembly:

.../organism/assembly/

- What we maintain:
 - Genome sequence (fasta)
 - Alignment indices: BWA, TopHat, Bowtie (1,2)
 - Annotations:
 - GTF format gene model
 - RefFlat format gene model
 - Signal artifact list (if available)

Working with Lustre

- I. Revisit architecture
- 2. Stripes
- 3. Avoiding I/O Bottlenecks
- 4. Using System Cache

Lustre: Quick Review

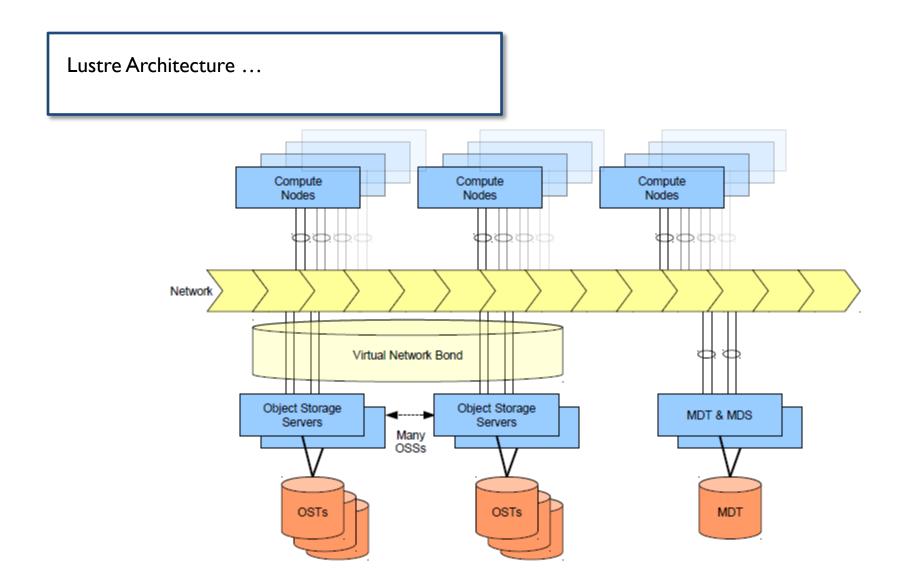
Lustre is a massively parallel distributed file system

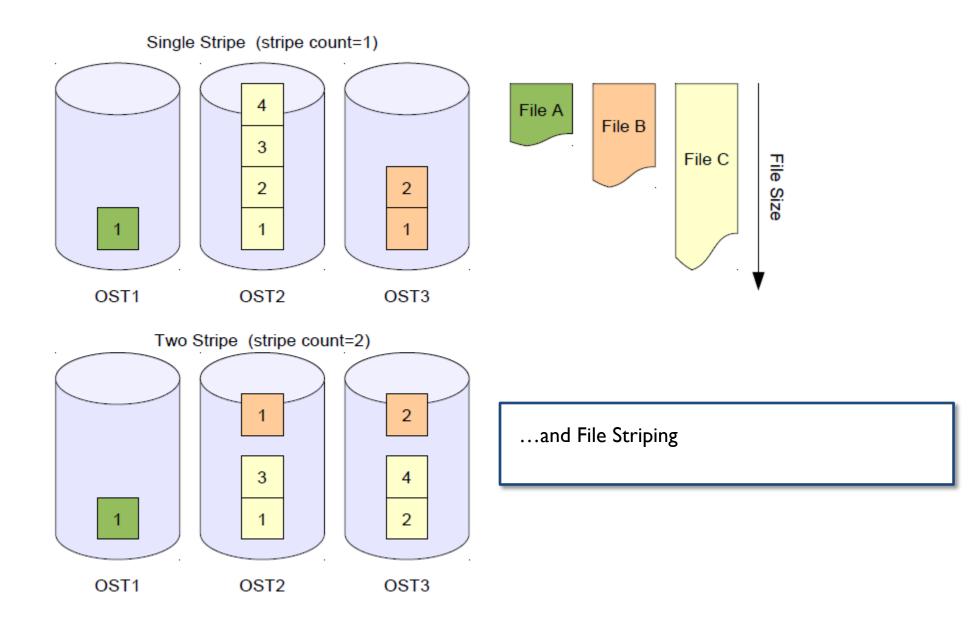
- Deployed in 7 out of 10 most powerful supercomputers
- POSIX compliant

Lustre design paradigm concepts

- Separation of file meta-data and storage allocation
- Scalable data serving through parallel data striping
- Aggregates network bandwidth
- Distributed operation

'**scratch**' storage we (deliberately) don't back it up





File Striping of Large Files

Performance benefits

- Aggregates I/O bandwidth to single large file
- In general, more stripes improves performance
- Small overhead associated with open/closing striped files

(Striping allows file size to exceed single OST size)

Many jobs reading single file

- For example blastdb and maq reference data
- /lustre/reference_data/genomes

Many jobs reading & writing multiple large files

- Requires benchmarking

Many jobs writing to single file

 High bandwidth but requires careful coding (can be disastrous)

Set Stripe Information

Set per file or directory Default is not to stripe Only newly created files will be stripes

- Use cp (not mv) to migrate existing files

clust1-headnode ~ \$ lfs setstripe <file|dir> --size <stripe_size>
--count <count> --index <index>

Where,

size = stripe size specified in k, m or g (0 default IMB)

```
count = OST stripe count (0 defaults 4 OST
```

```
and -I over all OSTs)
```

```
index = OST index of first stripe (-1 indicating default)
```

Read Stripe Information

Inspect file and directory stripe information with **lfs getstripe**

Using System Cache

Disk access is slow (no escape from this!)

- Memory access measured in a few nanoseconds
- Disc access measured in 10s of milliseconds

Linux uses free memory as cache

- Memory reclaimed as least used files expunged
- "Pre-warming" cache
- Can increase I/O performance

Clust1-headnode ~ \$ cat largefile > /dev/null Clust1-headnode ~ \$ grep searchString largefile

Avoiding Cache misses

Say you want to compare 3 sequences against 3 large databases:

"Out of the box" example

- sequence | vs database | No cache: disk read required
- sequence I vs database 2 Cache miss: disk read required
- sequence I vs database 3 Cache miss: disk read required
- sequence 2 vs database | No cache: data expunged from cache
- sequence 2 vs database 2 Cache miss: disk read required

Re-ordering to avoid cache misses

- sequence | vs database | No cache: disk read required
- sequence 2 vs database 1 Cache hit: data in cache
- sequence 3 vs database 1 Cache hit: data in cache
- sequence I vs database 2 No cache: disk read required
- sequence 2 vs database 2 Cache hit: data in cache

Avoiding Bad Performance

- \cdot Interactive use
- Statting files can be slow (common with shared file systems)
- Avoid directly editing small files on lustre (keep to /home)
- Turn off "color ls" (stat required for each file/directory)

Random seeks

- Small random I/O extremely slow on lustre
- Avoid, as much as possible, running databases on lustre
- e.g. mysql, sqlite, Berkeley DB etc

Limit number of files in directory

- 10,000s files in single directory bad (avoid, if possible)
- Use lfs setstripe to confine all files to single
 OST obviously for small files only!

Clust1-headnode ~ \$ lfs setstripe --count 1 directory

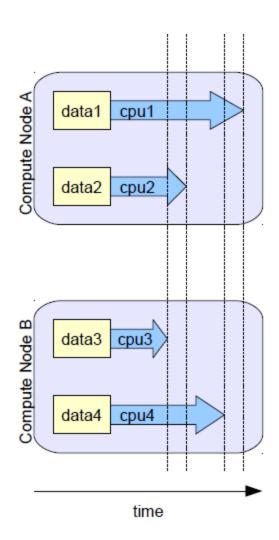
Simple Parallel

('Embarassing' or 'Trivial' in the computing science literature)

Solving many similar and independent tasks

- Analysis split into tasks
- Task assigned to one cpu
- No inter-task communication
- More throughput by running more tasks
- Task runtime varies

90% of bioinformatics codes fall into this model



Shared Memory

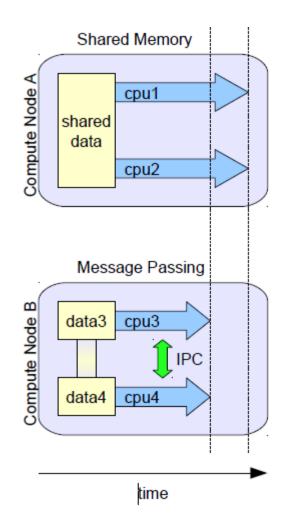
Shared tasks and memory

- Tasks assigned to cpus or cores
- Inter-task communication via shared memory
- Runtime decreases by adding more threads

Message Passing - local

Multiple processes communicate using O/S level systems. Code must be specifically written to exploit parallelism

- OpenMP/OpenMPI/etc



Message passing over network

Single task split across many compute nodes

- inter-machine communication (IMC)
- through MPI/OpenMP libraries again

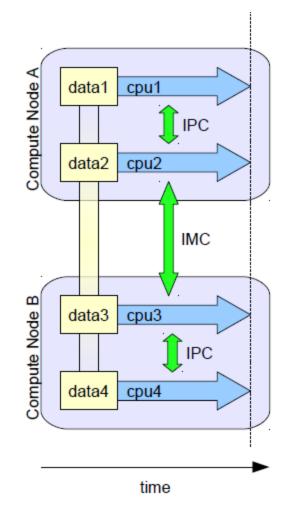
Hybrid models

Single task split across many compute nodes

- Mix SMP, local MP, network MP

Can be tricky to predict performance.

Your code may get quite complex...





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